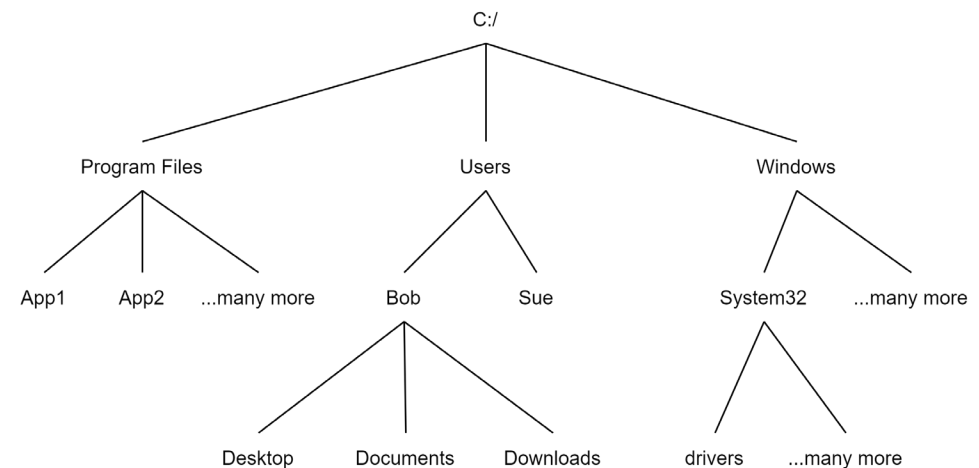
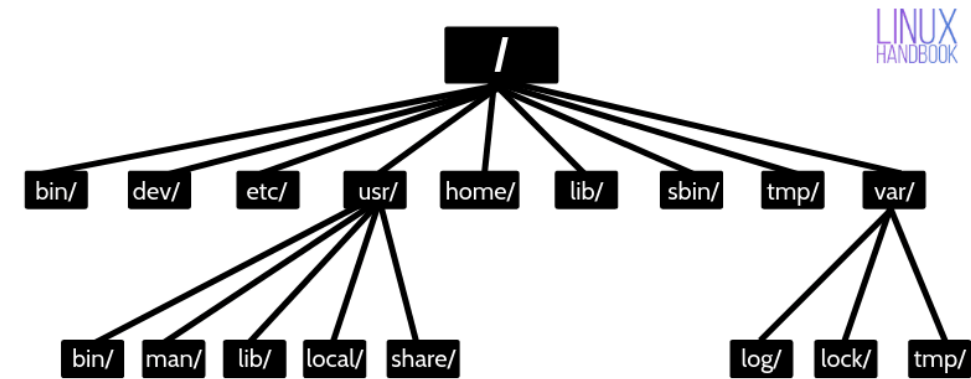
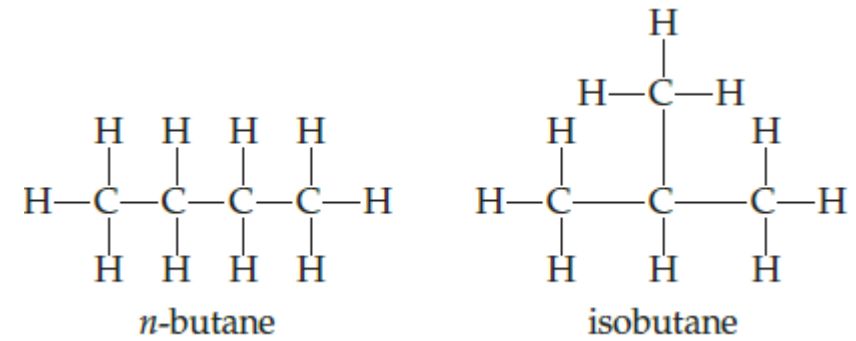


Trees



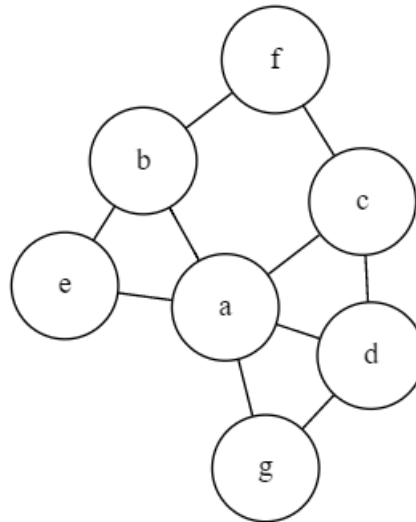
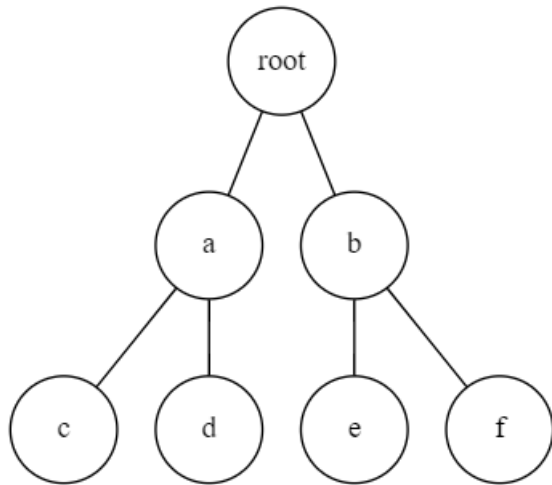
Background

- Arthur Cayley, English mathematician, discovered trees in 1857 when he needed a way to count chemical compounds
- Fast Searching
- Data compression
- Modeling games/procedures/decisions/relationships
- Filesystem hierarchy
- Evaluating mathematical expressions



Trees vs Graphs

- Both are composed of vertices and edges
- Tree: Only one path from one vertex to any other vertex
- Graph: More than one path between vertices



Tree Terminology

Free Tree

Rooted Tree

Root

Level

Height

Parent

Child

Ancestor

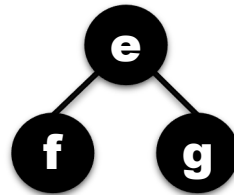
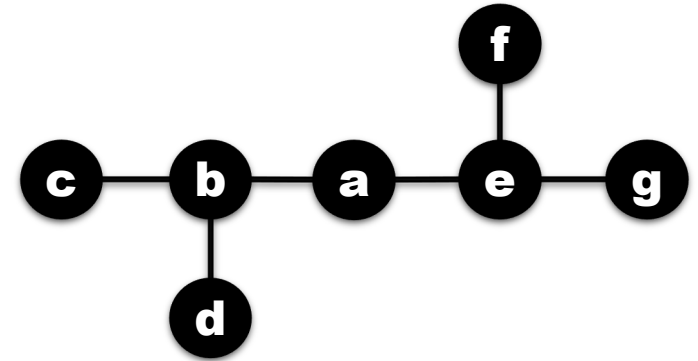
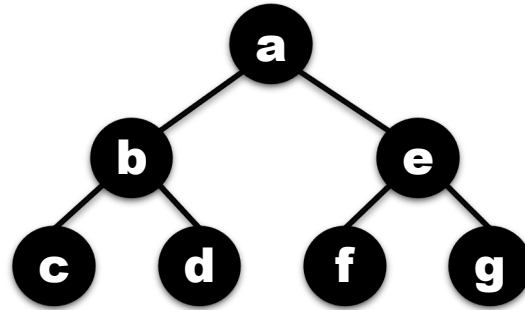
Descendant

Leaf

Internal vertex

Sibling

Subtree



Review:

What does it mean if two graphs are isomorphic?

With a partner, complete Additional Exercises:

14.1.2

A Few Tree Applications

Binary Search Trees

- Fast searching. Optimize **time**.

Huffman Trees

- Data compression. Optimize **space**.
- Binary tree, but not a binary **search** tree.
- Used to create an optimal prefix code

Game trees

Mathematical expressions

Consider a text file that consists of 1000 a's.

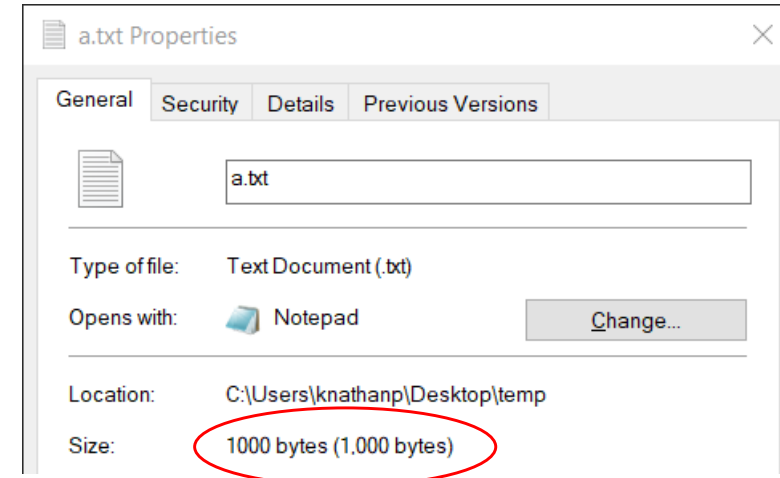
aaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaa
...

How much space would it take to store this file if we use standard ASCII or UTF-8, which each use 1-byte (8-bits) to store the letter "a"?

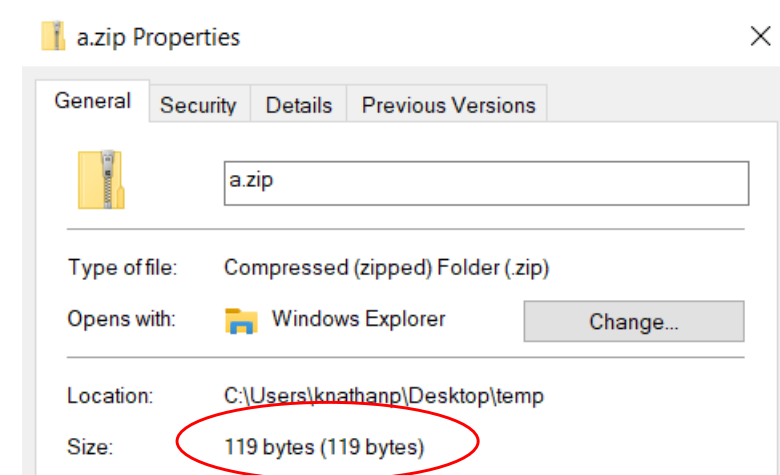
Can you think of a way to store this file in a way that would save space?

What is the compression ratio?

$$\frac{\text{original} - \text{compressed}}{\text{original}} = \frac{1000 - 119}{1000} = 0.88 = \mathbf{88\%}$$



ZIP compression

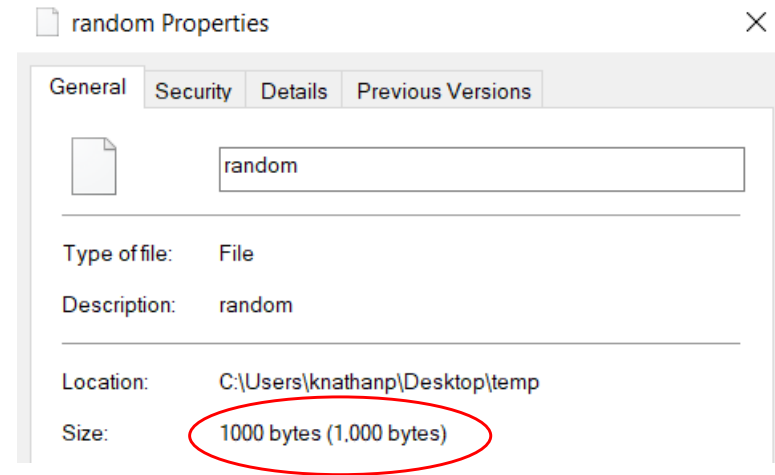


What about a file that consists of 1000 completely random bytes?

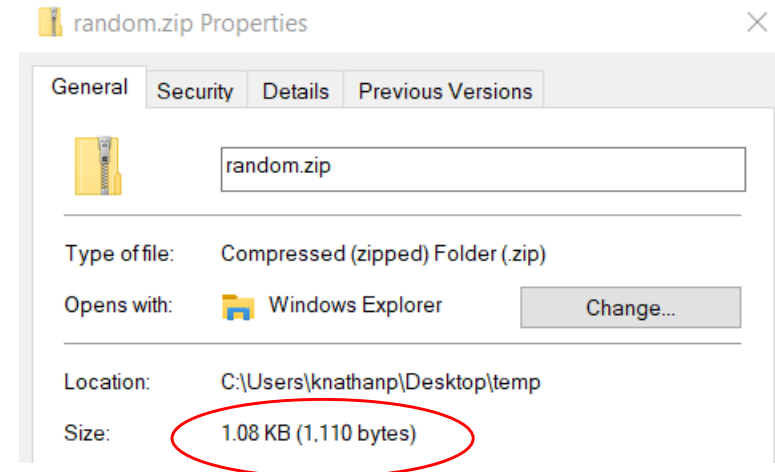
q2KVZRbqRiflgI56Es4plUeWoWBk
w2b892Ss5JWSJAK5SPaIksHucUo
GFMsZ2pqBB
...

Can you think of a way to store this file in a way that would save space?

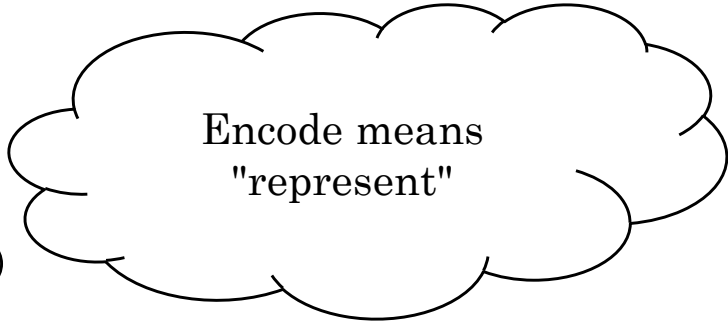
Why don't we save any space by compressing it?



ZIP compression



Fixed-length Encoding



Encode means
"represent"

- How many bits would it take to encode the letters of the English alphabet? (the minimum number of bits)
 - There are 26 letters
 - 4 bits (2^4) will only give 16 possible encodings
 - Need at least 5 bits
 - But that gives $2^5 = 32$ different possibilities. Waste of bits
 - There are six combinations of bits that are unused
- This is a **fixed-length** encoding. Each symbol/letter uses five bits.
- Wastes **space** and **time**.
- Can we find a coding scheme that uses fewer bits?

A	00000000	0 0
B	00000001	1 1
C	00000010	2 2
D	00000011	3 3
...	
Ø	11100001	14 25
Unused	11110010	26
Q	110011	27
R	11100	28
S	11101	29
T	11110	30
...	11111	31

Variable-length Encoding

- Instead of using a fixed number of bits to represent each letter/symbol, we can use bit strings of different lengths to encode letters.
- For efficiency, letters that occur more frequently should be represented using shorter bit strings. Letters that are rarely used should be represented using longer bit strings.
- Example:
The letters **e**, **a**, **t** are more common than x and z.

What if we represent the letter **e** with 0, **a** with 1, and **t** with 01?

Then the bit string 0101 could represent **eat**, **tea**, **eaea**, or **tt**.

e	0
a	1
t	01

Prefix Codes

- One way to make sure no bit string corresponds to more than one sequence of letters is to encode letters in such a way that the bit string for a letter **never occurs as the first part of a bit string for another letter**.
- This is called a Prefix Code.
- Example

To represent the letters **e**, **a**, **t** using a prefix code, we could choose 0 for **e**, 10 for **a**, and 11 for **t**.

The bitstring 10110 represents the word **ate**. There is no other possible interpretation.

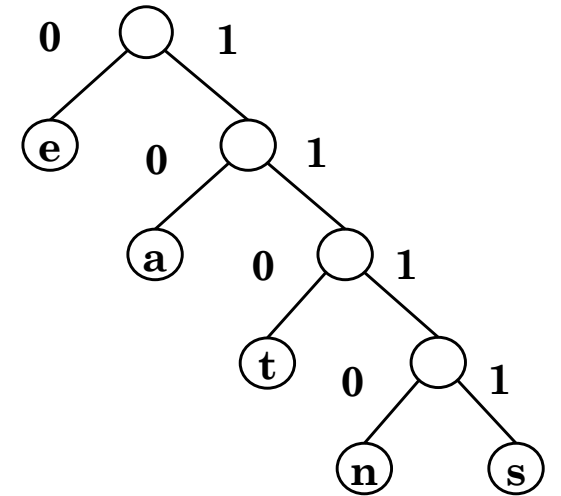
Now we can decode or recover a word from a bitstring without ambiguity.

e	0
a	10
t	11

Prefix Codes

- A prefix code can be represented using a binary tree.
- Leaves represent the characters/symbols we need to represent
- Edges represent the encoding bits. A left edge is a 0 and a right edge is a 1.
- The bit string used to encode a symbol is the sequence of edges to reach the symbol.
- Because this is a tree, the path to get to any leaf is **unique**.
- The sequence of bits formed by the path is a **prefix code**.

e	0
a	10
t	110
n	1110
s	1111



11111011100
S A N E

Prefix Code Examples

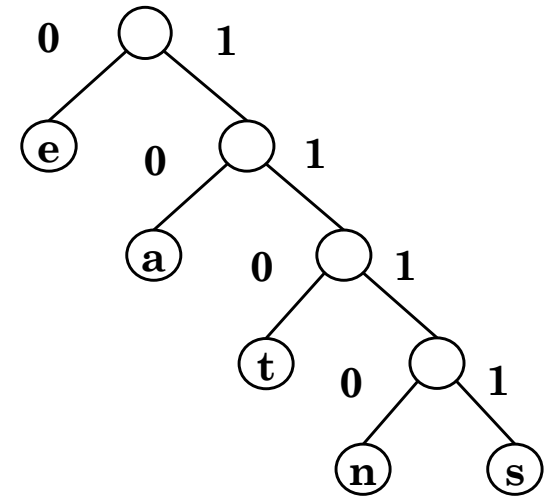
- Huffman codes
- [Country calling codes](#)
- ISBN country and publisher assignments
- Machine language instruction sets

Where is Huffman Coding used?

- [JPEG](#)
- [MP3](#)
- [DEFLATE algorithm \(Huffman Coding is combined with LZ77 compression\)](#)
 - zip, gzip, 7-Zip, C zlib library, PuTTY, ...
 - PNG files

Huffman Trees

- Used to produce a **Huffman Coding**, which is a prefix code.
- It is not only a prefix code, but it is a prefix code using the fewest number of bits possible.
- Fundamental algorithm in data compression.
- Uses a binary tree.
- Store more frequently seen symbols closer to root.
- Label the links rather than the nodes. Left link is 0, right link is 1.
- Symbols we want to represent are stored at the leaves.
- Each symbol is encoded into a unique bitstring by following the structure of the tree.



Compression Ratio

- How do we know how much space/time we are saving by using a variable-length encoding instead of a fixed-length encoding?
- Calculate the **compression ratio**!

Compression Ratio

f : bits per symbol for fixed length encoding

v : average bits per symbol with variable length encoding

$$\frac{f-v}{f} \cdot 100$$

Example: Create a Huffman Tree for the string "MISSISSIPPI"

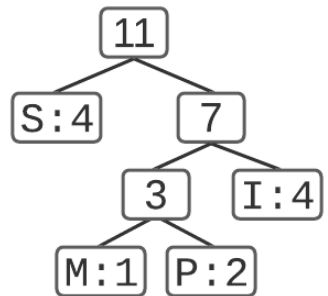
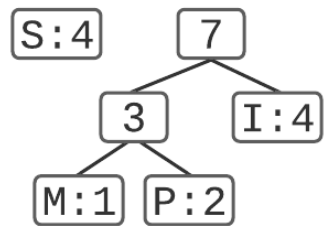
1. List out the counts

2. Go through the algorithm:

- 1. List symbols in a priority queue sorted by count (smallest count first)
- 2. Combine first two symbols (M and P) into a tree and add to the proper place in the queue
- 3. Repeat again, combining first two items in queue into a tree (M, P, and I) and adding back to the queue.
- 4. Continue to repeat until only one item is on the queue. This is the completed tree.

3. Extract the encoding for each letter:

100110011001110110111 21 bits
M I SSI SSI P P I



Letter	Count
M	1
P	2
I	4
S	4

Letter	encoding
M	100
P	101
I	11
S	0

Example: Create a Huffman Tree for the string "MISSISSIPPI"

Calculate the compression ratio.

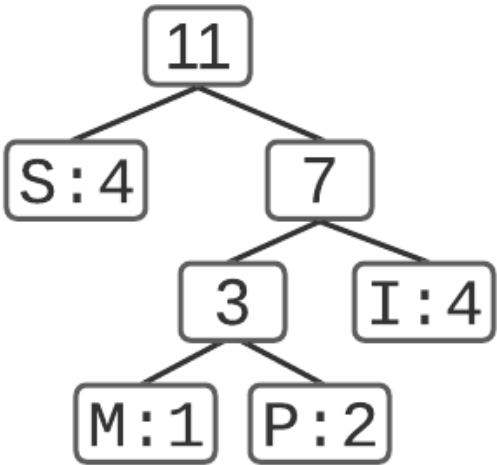
How many bits per letter would it take for a fixed-length encoding?

2 bits per letter

0010111110111110010110 22 bits

M I S S I S S I P P I

Letter	encoding
M	00
P	01
I	10
S	11



Letter	Count
M	1
P	2
I	4
S	4
total	11

How many bits on average with this Huffman encoding?

M: 3 bits * 1 letter = 3

P: 3 bits * 2 letters = 6

I: 2 bits * 4 letters = 8

S: 1 bits * 4 letters = 4

Total: 3+6+8+4 = 21

Average bits per letter: 21/11 = 1.91

Compression ratio:

$$\frac{f-v}{f} \cdot 100 \quad \begin{matrix} f = 2 \\ v = 1.91 \end{matrix}$$

$$\frac{2 - 1.91}{2} \cdot 100 = 4.5\%$$

Letter	encoding
M	100
P	101
I	11
S	0

Example: Create a Huffman Tree for the string "MISSISSIPPI"

0010111110111110010110

M I S S I S S I P P I

100110011001110110111

M I SSI SSI P P I

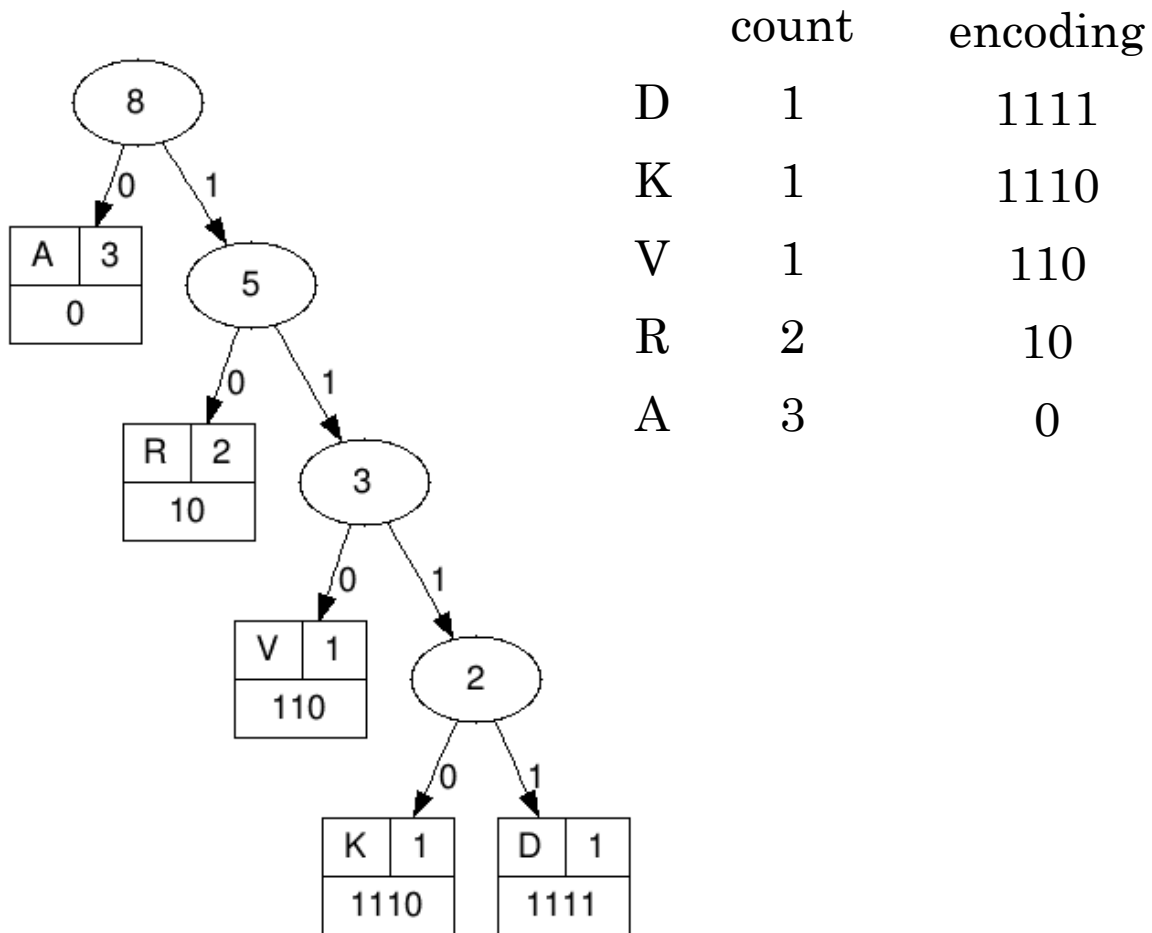
Fixed

Variable

Letter	fixed	variable
M	00	100
P	01	101
I	10	11
S	11	0

Try it:

Create a Huffman tree and Huffman encoding for "AARDVARK"



001011111100101110
AAR D V AR K

How many bits for fixed length encoding?

How many bits for variable length encoding?

Compression ratio?

Suggested Additional Exercises

14.2.2

14.2.3

Huffman Practical Application

[Python Example implementing Huffman Tree](#)