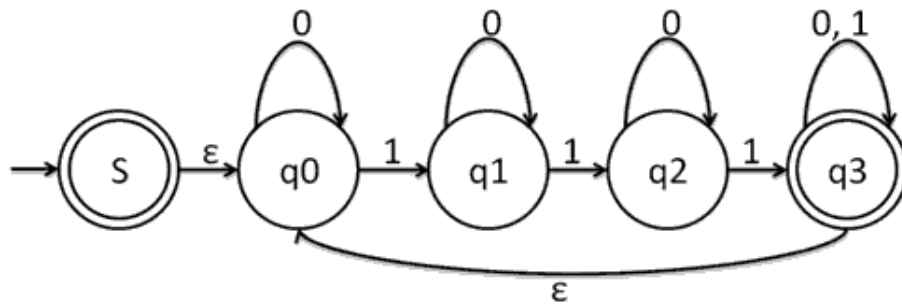


FSMs



Modeling Computation

- Modeling computation helps us to answer questions such as
 - Given a task, can it be done using a computer?
 - If so, how can the task be carried out?
- We can study computation using different structures
 - Grammars
 - Finite State Machines
 - Turing Machines
- In this class, we will touch briefly on Finite State Machines

Finite State Machines

Many types of machines, including computer components, can be modeled using **finite state machines**.

The basis for many types of programs:

- spell checkers
- grammar checkers
- searching text
- speech recognition
- markup language transformation (HTML, Markdown, LaTeX, etc.)
- network protocols
- ...

Finite State Machines

Different types of FSMs can model different things.

We will look at two types of FSMs:

1. Produce output
2. Recognize a language / bitstring / sequence of inputs / property

Finite State Machines

All FSMs include:

- Finite set of **states**
- **Starting state**
- **Input alphabet or input actions**
- **Transition function**

The FSM is given a sequence of inputs. The **transition function** determines the next state based on the current state and the current input.

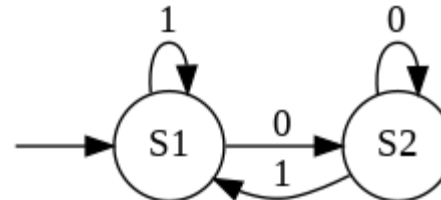
The transition function takes an ordered pair *(state, input)* and returns a *state*.

| | |
|------------------------------------|--------------------------------|
| Q | Set of states |
| $q_0 \in Q$ | Start state |
| I | Set of input actions / symbols |
| $\delta: Q \times I \rightarrow Q$ | Transition function |

Example:

$$\begin{aligned} Q &= \{S1, S2\} \\ q_0 &= S1 \\ I &= \{0, 1\} \\ \delta &= \{((S1, 0), S2), ((S1, 1), S1), ((S2, 0), S2), ((S2, 1), S1)\} \end{aligned}$$

We can also represent a transition function using a state diagram:



Finite State Machines

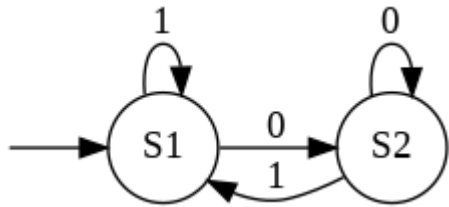
Example:

$$Q = \{S1, S2\}$$

$$q_0 = S1$$

$$I = \{0,1\}$$

$$\delta = \{((S1,0), S2), ((S1,1), S1), ((S2,0), S2), ((S2,1), S1)\}$$



Example Input: 1001101

Which is the last state given this input?

FSM with Output

Some FSMs produce **output** from an output alphabet or output set.

The **output** is determined based on the current state and the current input.

Example:

$Q = \{s0, s1, s2, s3\}$

$q_0 = s0$

$I = \{0,1\}$

$O = \{0,1\}$

$\delta = \{$

$(s0,0) \rightarrow (s1,1),$

$(s0,1) \rightarrow (s0,0),$

$(s1,0) \rightarrow (s3,1),$

$(s1,1) \rightarrow (s0,1),$

$(s2,0) \rightarrow (s1,0),$

$(s2,1) \rightarrow (s2,1),$

$(s3,0) \rightarrow (s2,0),$

$(s3,1) \rightarrow (s1,0)$

$\}$

How many ordered pairs are in the **domain** of the transition function δ ?

The domain is $Q \times I$. There are 4 states in Q and 2 input symbols, so $4 \times 2 = 8$

Go to the whiteboards and draw this FSM

What would be the output on the input of 1010101?

0111111

Input of 1100101?

0011010

Q

Set of states

$q_0 \in Q$

Start state

I

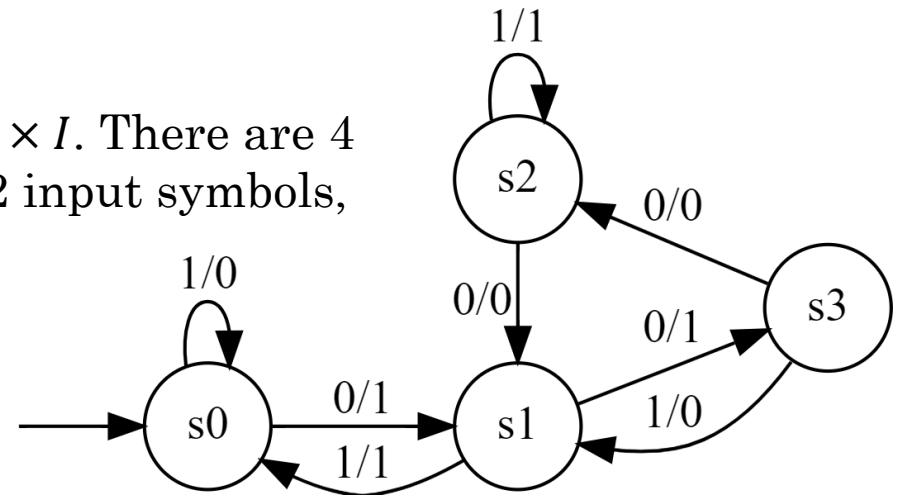
Set of input actions / symbols

O

Set of outputs

$\delta: Q \times I \rightarrow Q \times O$

Transition function

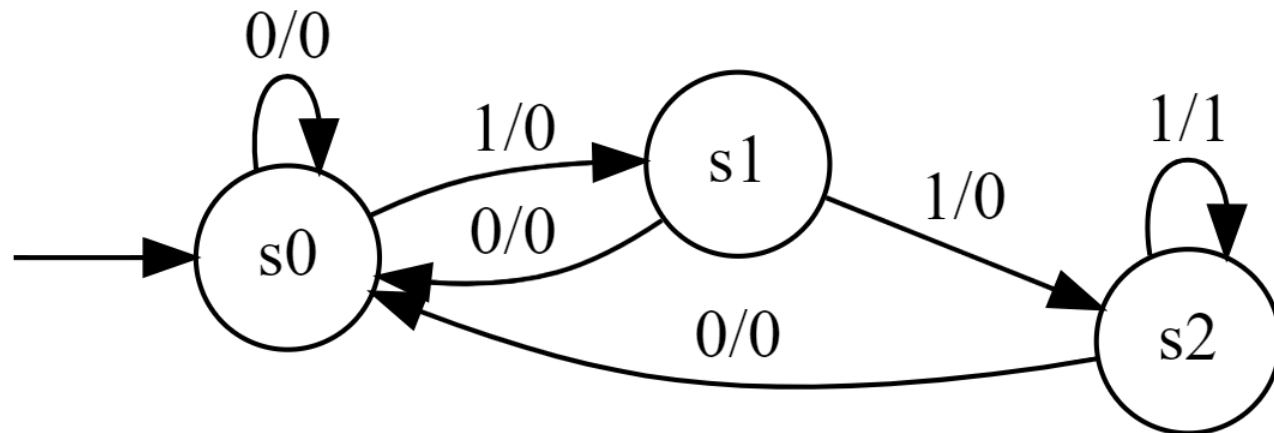


FSM with Output: Example

Imagine a coding scheme in which the receiver of a message knows there has been a transmission error whenever three consecutive 1s appear. Let's build an FSM that will output a 1 whenever it detects three consecutive 1s and a 0 otherwise.

An FSM can only remember its current state. So how many states will we need to "remember" three bits? **Three**

Let's draw it.



What would be the output on the input of 101101?

000000

On an input of 101110?

000010

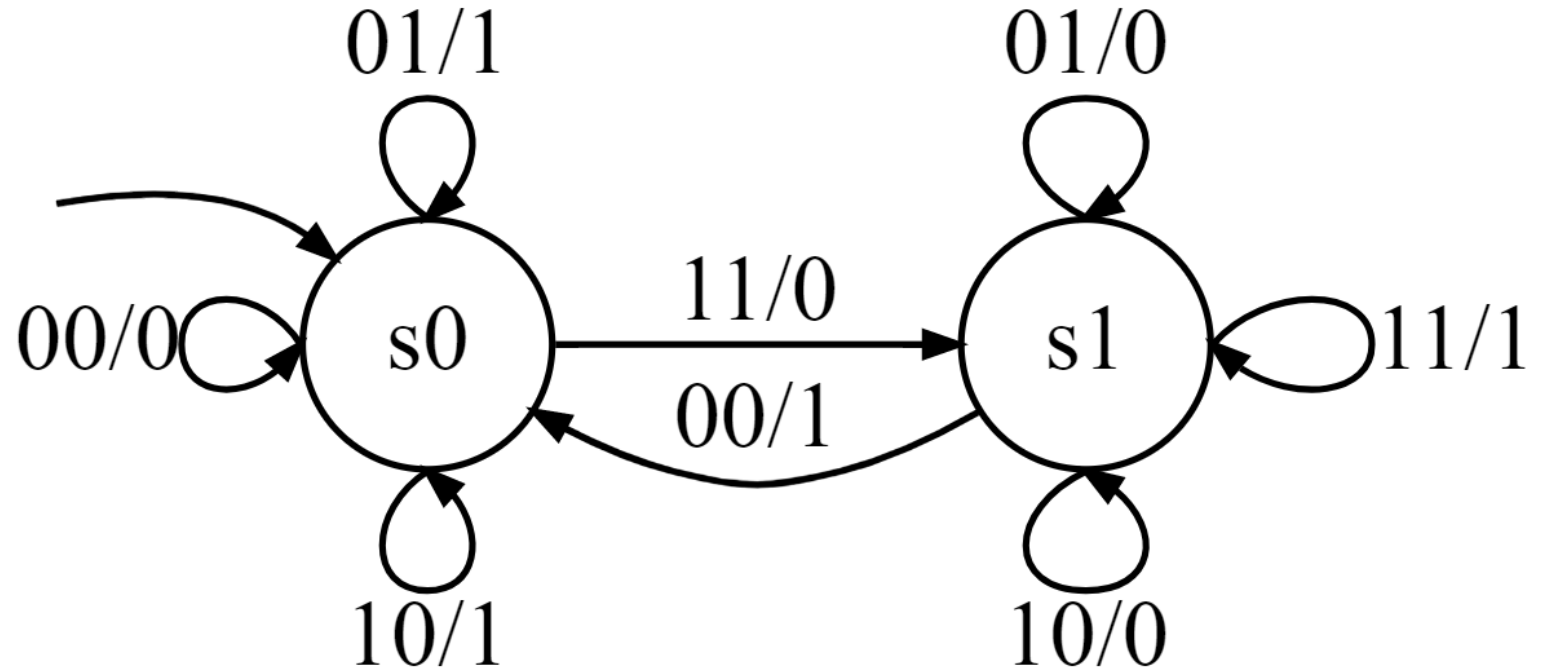


A "1" indicates the receiver detected an error

Example: Add two binary numbers

$Q = \{s0, s1\}$
 $q_0 = s0$
 $I = \{00, 01, 10, 11\}$
 $O = \{0, 1\}$
 $\delta = \{$
 $(s0, 00) \rightarrow (s0, 0),$
 $(s0, 01) \rightarrow (s0, 1),$
 $(s0, 10) \rightarrow (s0, 1),$
 $(s0, 11) \rightarrow (s1, 0),$
 $(s1, 00) \rightarrow (s0, 1),$
 $(s1, 01) \rightarrow (s1, 0),$
 $(s1, 10) \rightarrow (s1, 0),$
 $(s1, 11) \rightarrow (s1, 1)$
 $\}$

This FSM takes an input consisting of 2 bits at a time and outputs one bit. Now we can add two binary numbers.



What is the result of adding these two numbers by hand?

0110
0011

1001

Now trace these through the FSM, taking two bits at a time (one from each string from right to left). Does it produce the same result?

FSMs that Recognize Properties

Some FSMs recognize whether a sequence of inputs has a particular property.

Instead of producing an output, these FSMs have one or more **accept** states.

If an input string ends on an **accept** state, we say the input has the property recognized by the FSM, or that the input is a valid string in the language of the FSM.

Can be used to represent a **regular language** in a graphical way.

Example:

The language of all bitstrings that end with a 1:

"01", "11", "11111", "01010101", ...

Regular expression: $(0 \mid 1)^*1$

| | |
|------------------------------------|--------------------------------|
| Q | Set of states |
| $q_0 \in Q$ | Start state |
| I | Set of input actions / symbols |
| $A \subseteq Q$ | Accepting state(s) |
| $\delta: Q \times I \rightarrow Q$ | Transition function |



These types of FSMs are also known as **Finite State Automata**

Finite State Automata

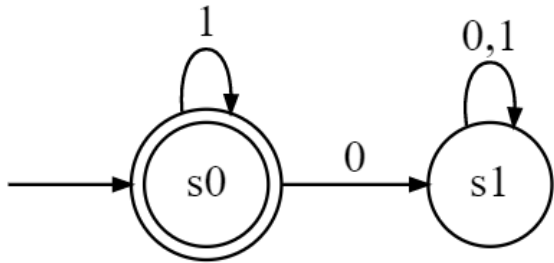
Do not produce output.

Have a set of final states (accepting states).

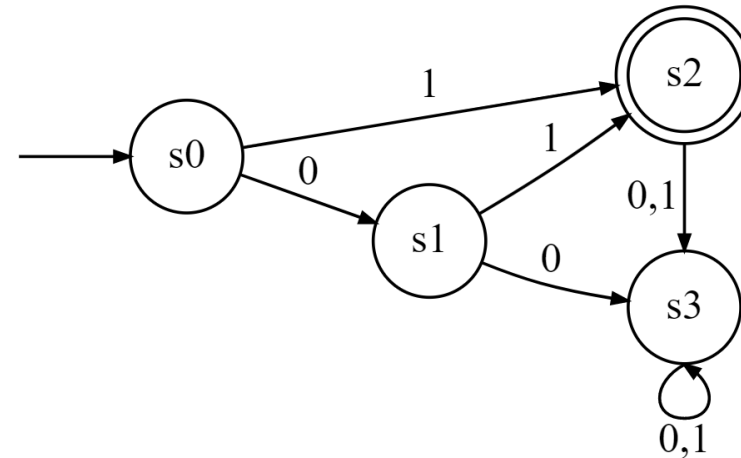
Accepting states are shown using a double circle.

| | |
|------------------------------------|--------------------------------|
| Q | Set of states |
| $q_0 \in Q$ | Start state |
| I | Set of input actions / symbols |
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Describe the language recognized by these finite state automata.



All bitstrings that contain only 1s

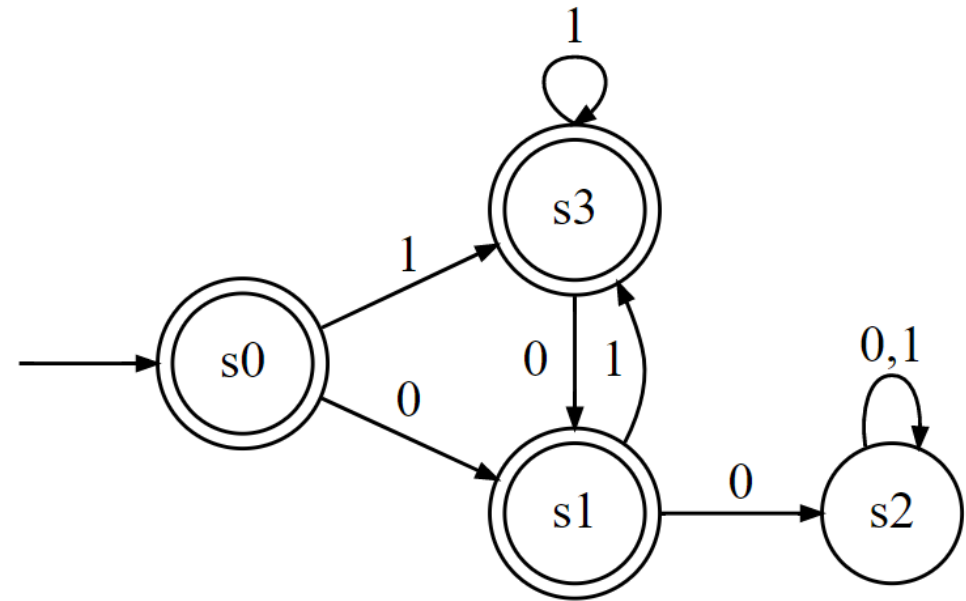
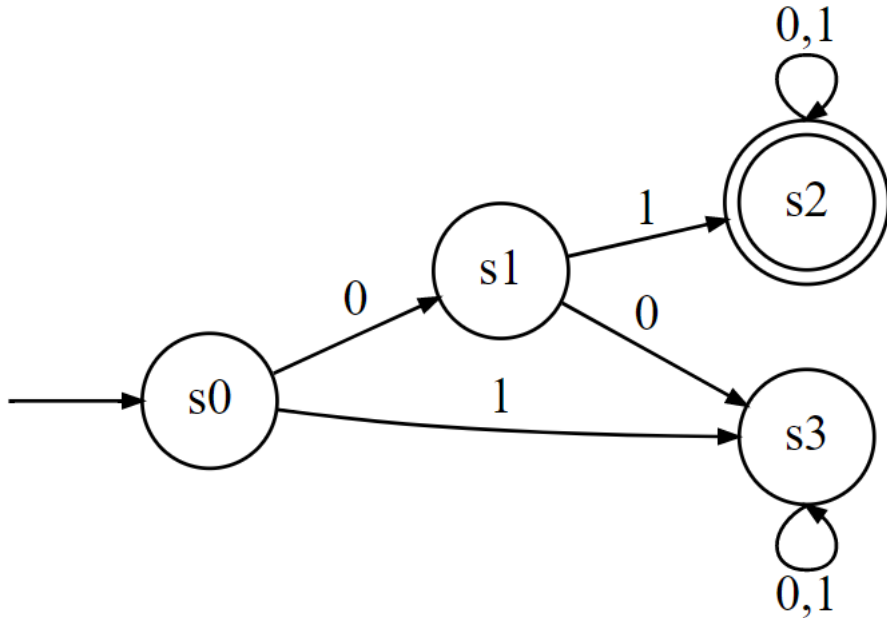


The bitstrings {1, 01}

Practice - Finite State Automata

Go to the whiteboards and draw these two finite state automata:

1. The set of bitstrings that begin with 01
2. The set of bitstrings that do not contain two consecutive 0s



FSM Examples in Python

[Examples of drawing FSMs using Python](#)